

**Amendments to the Claims**

1. (CURRENTLY AMENDED) A method of verifying the knowledge of a secret number  $s$  in a prover device(10) by a verifier device(30) having no knowledge of the secret number, with a zero-knowledge protocol using the Montgomery representation of numbers and Montgomery multiplication operations therein.

2. (ORIGINAL) The method of claim 1 in which the zero knowledge protocol is the Fiat-Shamir protocol.

3. (ORIGINAL) The method of claim 1 in which the zero knowledge protocol is the Guillou-Quisquater protocol.

4. (CURRENTLY AMENDED) The method of claim 2 including the steps of:

- (i) providing(102) to the verifier device(10) a value  $v = s^2$  being the Montgomery multiplication of the secret number  $s$  by itself;
- (ii) computing(106), by the prover device, the value  $x = r \times_m r$ , where  $r$  is a random number and transmitting the value of  $x$  to the verifier device;
- (iii) selecting(108), by the verifier device, a challenge value of  $e$  from the set  $\{0, 1\}$  and transmitting the challenge value to the prover device;
- (iv) computing(110), by the prover device, the value  $y = r \times_m s^e$  and transmitting the value  $y$  to the verifier device; and
- (v) the verifier device(10) checking the authenticity of the prover's response according to the values of  $x$ ,  $y$  and  $v$  previously received and according to the challenge value  $e$ .

5. (CURRENTLY AMENDED) The method of claim 4 wherein the step of checking the authenticity of the prover's response comprises the steps of:

for a challenge value of  $e = 1$ , computing(115) the values of  $y \times_m y$  and  $v \times_m x$  and checking(116) that they are the same; or

for a challenge value of  $e = 0$ , computing~~(120)~~ the value of  $y \times_m y$  and checking~~(121)~~ that it is the same as the previously received value of  $x$ .

6. (CURRENTLY AMENDED) The method of claim 4 ~~or claim 5~~ further including the steps of repeating steps (ii) to (v) for a number of consecutive rounds to confirm the authenticity of the prover device.

7. (CURRENTLY AMENDED) The method of claim 4 ~~or claim 5~~ in which the secret number  $s$  is a Montgomery representation of another number  $s'$  known in the prover device domain but not in the verifier device domain, further including the step of computing, by the prover device, the value of  $s$  from  $s'$  according to  $s = s'R \bmod n$ , where  $R > n$ , values of  $n$  and  $R$  being used by both the prover device and the verifier device.

8. (CURRENTLY AMENDED) The method of claim 4 in which the Montgomery multiplications of  $s \times_m s$ ,  $r \times_m r$ , and  $r \times_m s^e$  are carried out according to the formula  $a \times_m b = abR^{-1} \bmod n$ , where  $R > n$ , values of  $n$  and  $R$  being used by both the prover device~~(10)~~ and the verifier device~~(30)~~.

9. (CURRENTLY AMENDED) The method of claim 5 in which the Montgomery multiplications of  $y \times_m y$  and  $s^2 \times_m x$  are carried out according to the formula  $a \times_m b = abR^{-1} \bmod n$ , where  $R > n$ , values of  $n$  and  $R$  being used by both the prover device~~(10)~~ and the verifier device~~(30)~~.

10. (ORIGINAL) The method of claim 1 in which all computations in the zero knowledge protocol are performed using Montgomery representation of numbers and using Montgomery multiplication operations.

11. (CURRENTLY AMENDED) The method of claim 3 including the steps of:

(i) providing~~(303)~~ to the verifier device a value  $s^e$  being the Montgomery  $e^{th}$  power of the secret number  $s$ ;

(ii) computing ~~(306)~~, by the prover device ~~(10)~~, the value  $x = r^e$ , being the Montgomery  $e^{th}$  power of  $r$  where  $r$  is a random number, and transmitting the value of  $x$  to the verifier device ~~(30)~~;

(iii) selecting ~~(308)~~, by the verifier device, a challenge value of  $c$  from the set  $\{0, 1, \dots, e - 1\}$  and transmitting the challenge value to the prover device;

(iv) computing ~~(310)~~, by the prover device, the value  $y = r \times_m s^c$  and transmitting ~~(311)~~ the value  $y$  to the verifier device; and

(v) the verifier device ~~(30)~~ checking the authenticity of the prover's response according to the values of  $x$ ,  $y$  and  $s^e$  previously received according to the challenge value  $c$ .

12. (CURRENTLY AMENDED) The method of claim 11 wherein the step of checking the authenticity of the prover's response comprises the step of:

computing ~~(313, 314)~~ the values of  $y^e$  and  $x \times_m s^{ec}$  and checking that they are the same.

13. (CURRENTLY AMENDED) The method of claim 11 ~~or claim 12~~ further including the steps of repeating steps (ii) to (v) for a number of consecutive rounds to confirm the authenticity of the prover device.

14. (CURRENTLY AMENDED) A prover device ~~(10)~~ having contained therein a secret number  $s$  in Montgomery representation, the device adapted for proving the knowledge of the secret number  $s$  to a verifier device without conveying knowledge of the secret number itself, with a zero-knowledge protocol using the Montgomery representation of numbers and Montgomery multiplication operations therein.

15. (CURRENTLY AMENDED) The prover device of claim 14 further, including:

means ~~(12)~~ for selecting a random number,  $r$ ;

means ~~(11)~~ for computing the Montgomery square of  $r$  to obtain  $x$ ;

means for transmitting  $x$  to a verifier device ~~(30)~~;

means ~~(11)~~ for receiving a challenge value,  $e$ ;

means ~~(11)~~ for computing the Montgomery product of  $y = r \times_m s$ ; and  
means for transmitting  $y$  to the verifier device ~~(30)~~.

16. (CURRENTLY AMENDED) The prover device of claim 14 further including:

means ~~(12)~~ for selecting a random number,  $r$ ;  
means ~~(11)~~ for computing the Montgomery  $e^{th}$  power of  $r$  to obtain  $x$ ;  
means for transmitting  $x$  to a verifier device;  
means ~~(11)~~ for receiving a challenge value,  $c$ ;  
means ~~(11)~~ for computing the Montgomery product of  $y = r \times_m s$ ; and  
means for transmitting  $y$  to the verifier device.

17. (CURRENTLY AMENDED) A verifier device ~~(30)~~ for verifying the knowledge of a secret number  $s$  in a prover device without knowledge of the secret number itself, with a zero-knowledge protocol using the Montgomery representation of numbers and Montgomery multiplication operations therein.

18. (CURRENTLY AMENDED) The verifier device ~~(30)~~ of claim 17 further including:

means ~~(31)~~ for receiving the Montgomery square  $v$  of the secret number  $s$ ;  
means ~~(31)~~ for receiving the Montgomery square,  $x$  of a random number,  $r$ ;  
means ~~(31)~~ for transmitting a challenge value,  $e$  to the prover device;  
means ~~(31)~~ for checking the authenticity of the prover's response,  $y$  according to the Montgomery square of  $y$  verified against values of  $x$  and / or  $v$  received from the prover device according to the challenge value,  $e$ .

19. (CURRENTLY AMENDED) The verifier device of claim 17 further including:

means for receiving the Montgomery  $e^{th}$  power,  $s^e$  of the secret number  $s$ ;  
means for receiving the Montgomery  $e^{th}$  power,  $x$  of a random number,  $r$ ;  
means for transmitting a challenge value,  $c$  to the prover device;

means(31) for checking the authenticity of the prover's response,  $y$  according to the Montgomery  $e^{th}$  power of  $y$  verified against the value of  $x \times_m s^{ec}$  received from the prover device, according to the challenge value,  $c$ .

20. (CURRENTLY AMENDED) A method of proving the knowledge of a secret number  $s$  in a prover device(10) to a verifier device(30) having no knowledge of the secret number, with a zero-knowledge protocol using the Montgomery representation of numbers and Montgomery multiplication operations therein, comprising the steps of:

- selecting(305) a random number,  $r$ ;
- computing(306) the Montgomery  $e^{th}$  power of  $r$  to obtain  $x$ ;
- transmitting(306)  $x$  to a verifier device;
- receiving a challenge value,  $c$ ;
- computing(310) the Montgomery product of  $y = r \times_m s^c$ ; and
- transmitting(311)  $y$  to the verifier device.

21. (CURRENTLY AMENDED) A method of verifying the knowledge of a secret number  $s$  in a prover device(10) by a verifier device(30) having no knowledge of the secret number, with a zero-knowledge protocol using the Montgomery representation of numbers and Montgomery multiplication operations therein, comprising the steps of:

- receiving(103) the Montgomery square  $v$  of the secret number  $s$ ;
- receiving(107) the Montgomery square,  $x$  of a random number,  $r$ ;
- transmitting(108) a challenge value,  $e$  to the prover device;
- checking the authenticity of the prover's response,  $y$  according to the Montgomery square of  $y$  verified against values of  $x$  and / or  $v$  received from the prover device according to the challenge value  $e$ .

22. (CURRENTLY AMENDED) A method of verifying the knowledge of a secret number  $s$  in a prover device(10) by a verifier device(30) having no knowledge of the secret number, with a zero-knowledge protocol using the Montgomery representation of numbers and Montgomery multiplication operations therein, comprising the steps of:

receiving(303) the Montgomery  $e^{\text{th}}$  power of the secret number  $s$ ;  
receiving(307) the Montgomery  $e^{\text{th}}$  power,  $x$  of a random number,  $r$ ;  
transmitting(308) a challenge value,  $c$  to the prover device;  
checking(315) the authenticity of the prover's response,  $y$  according to the Montgomery  $e^{\text{th}}$  power of  $y$  verified against the value of  $x \times_m s^{ec}$  received from the prover device according to the challenge value  $c$ .

23. (CURRENTLY AMENDED) A computer program product, comprising a computer readable medium having thereon computer program code means adapted, when said program is loaded onto a computer, to make the computer execute the procedure of ~~any one of claims 1 to 13 and 19 to 22~~claim 1.

24. A computer program, distributable by electronic data transmission, comprising computer program code means adapted, when said program is loaded onto a computer, to make the computer execute the procedure of ~~any one of claims 1 to 13 and 19 to 22~~claim 1.

25. (ORIGINAL) Apparatus substantially as described herein with reference to the accompanying drawings.

26. (ORIGINAL) A method substantially as described herein with reference to the accompanying drawings.